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Published in:
I E E Transactions on Wireless Communications

DOI (link to publication from Publisher):
[10.1109/TWC.2013.032113.120975](https://doi.org/10.1109/TWC.2013.032113.120975)

Publication date:
2013

Document Version
Early version, also known as pre-print

[Link to publication from Aalborg University](#)

Citation for published version (APA):
Fyhn, K., Jensen, T. L., Larsen, T., & Jensen, S. H. (2013). Compressive Sensing for Spread Spectrum Receivers. *I E E Transactions on Wireless Communications*, 12(5), 2334 - 2343.
<https://doi.org/10.1109/TWC.2013.032113.120975>

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Compressive Sensing for Spread Spectrum Receivers

Karsten Fyhn, *Member, IEEE*, Tobias L. Jensen, *Member, IEEE*, Torben Larsen, *Senior Member, IEEE*, and Søren Holdt Jensen, *Senior Member, IEEE*

Abstract—With the advent of ubiquitous computing there are two design parameters of wireless communication devices that become very important: power efficiency and production cost. Compressive sensing enables the receiver in such devices to sample below the Shannon-Nyquist sampling rate, which may lead to a decrease in the two design parameters. This paper investigates the use of Compressive Sensing (CS) in a general Code Division Multiple Access (CDMA) receiver. We show that when using spread spectrum codes in the signal domain, the CS measurement matrix may be simplified. This measurement scheme, named Compressive Spread Spectrum (CSS), allows for a simple, effective receiver design. Furthermore, we numerically evaluate the proposed receiver in terms of bit error rate under different signal to noise ratio conditions and compare it with other receiver structures. These numerical experiments show that though the bit error rate performance is degraded by the subsampling in the CS-enabled receivers, this may be remedied by including quantization in the receiver model. We also study the computational complexity of the proposed receiver design under different sparsity and measurement ratios. Our work shows that it is possible to subsample a CDMA signal using CSS and that in one example the CSS receiver outperforms the classical receiver.

Index Terms—Compressive sensing, sparse sampling, spread spectrum receivers, multiuser decoding

I. INTRODUCTION

As wireless communication devices are becoming more and more widespread and ubiquitous, the need for power efficiency and low production cost becomes paramount. A power costly operation in wireless communication is the conversion from analog to digital signals - the Analog to Digital Converter (ADC). The classic ADC uses the Shannon-Nyquist sampling theorem to represent an analog signal in digital form. The Shannon-Nyquist sampling theorem states that to perfectly represent an analog signal, it must be sampled at a frequency higher than twice the signal's bandwidth. When this theorem is obeyed, the original analog signal may be reconstructed perfectly from its sampled representation. The Shannon-Nyquist sampling theorem has been the foundation of digital signal processing for more than half a century and is considered a fundamental building block of digital signal processing systems. Recently, a new concept termed Compressive Sensing (CS) [1], [2] has been attracting more and more attention in the signal processing community as it provides an exception to the lower bound on the sampling rate by exploiting sparsity

in the signal. If a signal is sparse in some arbitrary basis, it may be sampled at a rate lower than the Nyquist frequency. Sparsity in CS is when a signal is comprised of only a few atoms from a given basis. The sampled signal must be acquired through some linear measurement scheme. Examples of these are random Gaussian, Bernoulli and Rademacher measurement schemes, as well as the Random Demodulator (RD) [3], [4] and the Modulated Wideband Converter [5].

Compressive sensing has primarily been studied in the general signal processing area, and relatively few researchers have looked into its application in communication systems. In [6], [7] the authors examine the use of CS in Ultra-Wideband (UWB) communication systems for channel estimation where the sparsity of the signal lies in the time domain. Others have used compressive sensing for source coding in communication networks, together with network coding [8]. In the spread spectrum area, some researchers have studied the general use of CS for spread spectrum communication systems [9]. However, their work is mainly focused on using CS for fast multi-user detection, rather than subsampling. Another example is in [10], where the authors use CS to decrease the sampling rate of a GPS receiver by exploiting sparsity in the number of possible signal components at the receiver. Their receiver structure is based on possibly complicated hardware filters, which may make their implementation very difficult considering the impact of hardware filters to CS performance [11]. In [12] the authors treat a similar topic where they design spread spectrum codes to enable a base station to perform multi-user detection on a large number of users, of which only a few are active at a time. Their work focuses on simple on-off signalling, i.e. the existence of nodes, rather than communication with them, and solves the multi-user detection problem using an adapted convex optimization algorithm. Their motivation is on increasing the number of active users in a network, rather than decreasing the sampling rate of the ADC. A more applied approach is taken in [13] where compressive signal processing [14] is applied to enable subsampling of an IEEE 802.15.4 Direct Sequence Spread Spectrum (DSSS) signal. In [15] the authors also solve a multiuser detection problem using compressive sensing, but in their work the focus is on the design of possibly complex analog filters. For this paper we focus on keeping the analog part as simple as possible and process the signals in the digital domain instead.

In our work we apply CS to a general CDMA system. We show that a RD implementation may be used to subsample the CDMA signal, but we also develop a simplified version of

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the RD which performs equally well for CDMA signals but is simpler and cheaper to implement. Our motivation is that by taking fewer samples we may be able to conserve power in the receiver, as can be seen in e.g. Eqn. 13 in [16]. We show the performance of the proposed receiver structure for the simple discrete case, when compared to a classic receiver structure and an RD receiver structure. Then we extend our results to a full RF numerical simulation and demonstrate that the performance is identical in this setting. Due to noise folding the CS approach suffers a penalty for downsampling, but we show that if quantization is taken into account CS outperforms the classic receiver in some cases. We finally investigate the complexity of the developed algorithms and compare the computational cost of the numerical experiments with the theoretically calculated computation cost. Following the paradigm of Reproducible Research [17], all our results and code are made available at <http://www.sparsesampling.com/css>.

To define our notation, let all vectors and matrices be denoted using lower and upper case letters in bold, \mathbf{x} and \mathbf{X} , respectively. The Penrose-Moore pseudo-inverse is denoted as \mathbf{X}^\dagger , the transpose of a real matrix as \mathbf{X}^T and the conjugate transpose of a matrix as \mathbf{X}^* . The expectation operator is denoted by $\mathbb{E}[\cdot]$.

In the following, we first develop a simple signal model in Section II, based on a dictionary of Gold sequences. We then elaborate on what CS is and which reconstruction algorithm we use in the numerical experiments in Section III. Furthermore, we define a novel measurement matrix design for spread spectrum receivers and demonstrate numerically how this measurement matrix performs with a Gold dictionary and the Subspace Pursuit reconstruction algorithm. This performance is compared to that of a Rademacher measurement matrix and a RD measurement matrix. This is followed by Section IV, which includes a simple numerical experiment of the different receiver structures. In Section V we extend the experiment to a full RF simulation with and without quantization. We then analyze the computational complexity of the proposed method in Section VI, after which we conclude the paper in Section VII.

II. SIGNAL MODEL

First, we consider a purely discrete model of a spread spectrum communication system. Uncoded information bits are sent in a slotted fashion, with each slot containing a single CDMA signal. The system is assumed to be synchronized, which may be obtained by e.g. having a central node or base station transmit beacons, which signify the beginning and end of slots. This is how mobile phone networks and some wireless sensor networks operate. The receiver is considered non-coherent, as information is encoded in the phase, but we do assume that there is no carrier frequency offset between the transmitter and receiver oscillators. This is of course not a realistic assumption but it is done to keep the system simple. Future work should investigate the impact of oscillator drift on performance. Each slot contains an independent CDMA signal and the slots are decoded sequentially and independently of each other.

For one slot, define a discrete QPSK baseband signal, $\mathbf{x} \in \mathbb{C}^{N \times 1}$ as:

$$\mathbf{x} = \Psi \boldsymbol{\alpha}, \quad (1)$$

where $\Psi \in \mathcal{S}_\Psi \subset \{\pm 1\}^{N \times N}$ is an orthogonal or near-orthogonal dictionary, containing spreading waveforms for transmission, \mathcal{S}_Ψ is the subset of $\{\pm 1\}^{N \times N}$ that contains orthogonal or near-orthogonal dictionaries and $\boldsymbol{\alpha} \in \{\pm 1 \pm j, 0\}^{N \times 1}$ is a sparse vector, that selects which spreading waveform(s) and what QPSK constellation point(s) to send. $\boldsymbol{\alpha}$ is a vector here because we only process one slot at a time and we assume that within a slot, the signal amplitude for each user is constant. That $\boldsymbol{\alpha}$ is assumed to be sparse is justified in some scenarios, which is demonstrated shortly.

An example of a system using the above signal model could be a wireless sensor network in which one node must gather information from any possible neighbors. Each node has a unique CDMA sequence assigned, which it uses to transfer information and each node does not know which neighbors it has, but it knows all possible CDMA sequences. Note that in this signal model $\boldsymbol{\alpha}$ is defined so that all users have identical amplitude. This is not realistic as the distance between nodes might vary a lot, resulting in differences between amplitude in the received signal components. We choose this simplification here but the reconstruction algorithm is not limited by this and also works for sparse vectors with different amplitude components.

In cases where the number of active nodes or users in a network is smaller than the total number of possible users, the vector $\boldsymbol{\alpha}$ may be assumed sparse, which is the enabling factor for CS. Cases such as these arise in e.g. mobile phone networks and wireless sensor networks, where the number of surrounding nodes *may* be large, but is *often* small.

At the receiver the following signal is observed:

$$\mathbf{y} = \Theta (\mathbf{x} + \mathbf{w}) = \Theta \Psi \boldsymbol{\alpha} + \Theta \mathbf{w}, \quad (2)$$

where Θ is a measurement matrix, which we shall treat later, and $\mathbf{w} \in \mathbb{C}^{N \times 1}$ is Additive White Gaussian Noise (AWGN). Notice here that we take into account noise folding as the noise is folded down into the compressed domain together with the signal. This makes the noise colored and has an impact on the demodulation performance, because each time the sampling rate is reduced by one half, the Signal to Noise Ratio (SNR) is decreased by 3 dB [18], [19].

A. Spread Spectrum Dictionary of Gold Sequences

In spread spectrum signals, a possible dictionary Ψ is a set of Gold sequences, as used in e.g. GPS technology [20]. A set of Gold sequences is a special dictionary of binary sequences with very low auto and cross-correlation properties [21]. To generate a Gold dictionary, two maximum length sequences must be generated by two linear feedback shift registers (LFSR). A maximum length sequence is often denoted an m -sequence (it has m elements), and is a special kind of pseudo-random noise sequence generated by a LFSR, such that it is periodic and produces a sequence of length $2^m - 1$. It is called a maximum length sequence as its period is at maximum

length. The reason for the length being $2^m - 1$ rather than 2^m is that the state where all cells are zero must be avoided. To obtain an m -sequence, the LFSR must be carefully chosen as there is no algorithm for ensuring maximum length. However, there are many known LFSR setups for varying choices of m . Furthermore, the two m sequences must be chosen so that their periodic cross-correlation is three-valued and takes on only the values $\{-1, -t, t-2\}$, where:

$$t = \begin{cases} 2^{(m+1)/2} + 1 & \text{for odd } m \\ 2^{(m+2)/2} + 1 & \text{for even } m. \end{cases} \quad (3)$$

The set of Gold sequences are then generated using two m -sequences: \mathbf{g}_1 and \mathbf{g}_2 , both of length $N = 2^m - 1$. Each Gold sequence in the set is generated as $\mathbf{g}_1 \oplus \mathbf{g}_i$ (exclusive or), where \mathbf{g}_i is \mathbf{g}_2 cyclically shifted by the parameter i . As i can take on values between $1 \leq i \leq 2^m - 1$, each shift constitutes a candidate for the set, resulting in a dictionary as follows: Define a $N \times N$ dictionary of Gold sequences as Ψ , with each column signifying a possible code sequence.

When using such a CDMA dictionary, the received signal must be sampled at a rate corresponding to the chip rate, where a chip is one entry in the received Gold sequences. If α is sparse the information rate of the signal is much lower and it may be possible to decrease the sampling rate by using CS.

In this paper, we use three Gold dictionary sizes: $m = 5, m = 7$ and $m = 10$. The m -sequence feedback sets used to generate these are:

- $m = 5$: $X^5 + X^2 + 1$ and $X^5 + X^4 + X^3 + X^2 + 1$
- $m = 7$: $X^7 + X^6 + 1$ and $X^7 + X^4 + 1$
- $m = 10$: $X^{10} + X^3 + 1$ and $X^{10} + X^9 + X^8 + X^6 + X^3 + X^2 + 1$

The chosen polynomials may be validated by calculating the auto and cross-correlation of the generated dictionaries and verifying that they follow the structure listed in the above.

III. COMPRESSIVE SENSING

CS is a novel sampling scheme, developed to lower the number of samples required to obtain some desired signal. At the heart of CS is the linear sampling scheme, called the measurement matrix. In classic receivers the measurement matrix Θ may be modelled as the identity matrix, such that \mathbf{x} is sampled at the chip rate of each channel (I and Q). Here, we shall denote a classic receiver using $\Theta_1 = \mathbf{I}$, where the subscript 1 denotes no subsampling and \mathbf{I} is the identity matrix of size $N \times N$. In CS another measurement matrix is used. Denote by $\Theta_\kappa \in \mathbb{R}^{M \times N}$ a CS measurement matrix, where $\kappa \in \mathbb{N}_1$ is the subsampling ratio when compared to the Nyquist rate and $M = N/\kappa$ (If κ does not divide N , M is rounded to the nearest integer). This measurement matrix is then responsible for mapping the N -dimensional signal \mathbf{x} to a M -dimensional signal \mathbf{y} . Normally this would make it impossible to recover the original signal, but under the assumption that \mathbf{x} is sparse in some basis, it is possible to reconstruct the original signal from the sampled, M -dimensional signal \mathbf{y} [1], [2].

Notice that we are not interested in the reconstructed signal, \mathbf{x} , but in the sparse vector α , which allows us to demodulate the data in the signal. We may obtain an estimate of α by reconstructing the signal from \mathbf{y} . Such a reconstruction may be obtained using e.g. a convex optimization problem solver

or a greedy algorithm. For this work, we choose the greedy algorithm Subspace Pursuit [22]. This algorithm is chosen due to its good performance in terms of both reconstruction accuracy and running time, as shown in Section III-B.

Before explaining the reconstruction algorithm, we return to the measurement matrix and introduce a new measurement scheme which is enabled by the use of CDMA. This new measurement scheme is easier to implement than the RD, but performs almost identically for spread spectrum systems. We call this a Compressive Spread Spectrum (CSS) measurement matrix and explain it further in the following.

A. Compressive Spread Spectrum Measurement Matrix

In most CS literature a choice of measurement matrix or structure must be made. The Bernoulli or Rademacher distributed measurement matrix is often seen in the theoretical literature, but it is not well suited for practical implementation in a wireless receiver. The Random Demodulator (RD) sampling structure [3], [4] is one of the most well-known measurement matrix structures developed, which is well suited for practical implementation. In the RD a Pseudo-Random Noise (PRN) sequence is mixed with the received signal followed by low-pass filtering. Because a spread spectrum transmitter has already spread the signal before transmission, we show that the RD structure can be improved so that the mixing with a PRN sequence at the receiver may be skipped. This is similar to what is done in [13] with IEEE 802.15.4 signals, which uses Direct-Sequence-Spread-Spectrum (DSSS) signals. These can be viewed as a special class of CDMA signals, which are used to counter interference from blockers in the same frequency band, rather than to distinguish between users or signals.

The proposed measurement matrix may therefore be defined similarly to the definition of the RD matrix in [4]. In their work, the measurement matrix is based on two matrices, \mathbf{D} and \mathbf{H} . First, let $\epsilon_0, \epsilon_1, \dots, \epsilon_N \in \{\pm 1\}$ be the chipping sequence used in the RD for a signal of length N . The mapping $\mathbf{x} \rightarrow \mathbf{D}\mathbf{x}$ signifies the demodulation mapping with the chipping sequence, where \mathbf{D} is the diagonal matrix:

$$\mathbf{D} = \begin{bmatrix} \epsilon_0 & & & \\ & \epsilon_1 & & \\ & & \ddots & \\ & & & \epsilon_N \end{bmatrix}. \quad (4)$$

Second, the \mathbf{H} matrix denotes the accumulate-and-dump action performed after mixing. Let M denote the number of samples taken and assume here that M divides N . Then each sample is the sum of N/M consecutive entries of the demodulated signal. The matrix performing this sampling action may therefore be defined as an $M \times N$ matrix, with N/M consecutive unit entries in the r th row starting in column $rN/M + 1$ for each $r = 0, 1, \dots, M - 1$. An example with $M = 3$ and $N = 6$ is:

$$\mathbf{H} = \begin{bmatrix} 1 & 1 & & & & \\ & & 1 & 1 & & \\ & & & & 1 & 1 \end{bmatrix}. \quad (5)$$

The RD is therefore designed to sample an analog signal, so that in a discrete representation this is the equivalent to:

$$\mathbf{y} = \mathbf{H}\mathbf{D}\mathbf{x}, \quad (6)$$

where \mathbf{x} is the Nyquist sampled input signal and \mathbf{y} is the compressively sampled output signal.

The reason for applying a chipping sequence is to spread the signal across the frequency spectrum, so that information is aliased down into the lower frequency area, which is left untouched by the low-pass filtering. In the proposed receiver this mixing is unnecessary because the signal has already been spread at the transmitter. The proposed receiver may therefore be simplified to:

$$\mathbf{y} = \mathbf{H}\mathbf{x}. \quad (7)$$

This is significantly simpler to implement in hardware than the RD. Comparing to the notation introduced for the measurement matrix in Section II we therefore have: $\Theta_\kappa = \mathbf{H}$.

To justify the use of no PRN sequence in the measurement matrix, consider the following. The use of a CDMA dictionary introduces a random-like dictionary matrix, which spreads the signal out so that each sample contains a little bit of the original information signal. This is similar to what the measurement matrix does in CS. Therefore, the sampling process may be rewritten as:

$$\mathbf{y} = \mathbf{H}\mathbf{x} = \mathbf{H}\Psi\alpha = \Theta\mathbf{I}\alpha. \quad (8)$$

Here, the measurement matrix becomes $\Theta = \mathbf{H}\Psi$, i.e. the subsampling matrix and the CDMA codes. The dictionary then becomes the identity matrix. When viewed like this, it is clear that Θ and \mathbf{I} are incoherent as the identity matrix only takes out one element in Θ . Another common mathematical tool for verifying the validity of a measurement scheme for compressive sensing is Restricted-Isometry-Property (RIP). However, the RIP gives a less precise and more conservative boundary between reconstruction success and failure than other bounds, see e.g. the discussions in [23], [24]. Instead, phase-transition diagrams [23] may be used to demonstrate empirically for which levels of sparsity the dictionary and measurement matrix are applicable. In the following, we first define the Subspace Pursuit algorithm and then use phase-transition diagrams to show that the proposed CSS measurement matrix has transitions that are very close to those of the Rademacher and RD measurement matrices for dictionary matrices using Gold sequences.

B. Subspace Pursuit

To reconstruct the signal a reconstruction algorithm must be chosen. Many different approaches have been developed, but two main classes of reconstruction algorithms are in widespread use: ℓ_1 minimization and greedy algorithms. Often, ℓ_1 minimization provides the best solution, but if the matrices Ψ and Θ are very large, it is much more efficient to use the simpler greedy algorithms [25]. Therefore, we choose to use greedy algorithms in this work.

In [25] an extensive numerical comparison between reconstruction algorithms is performed based on phase transition

plots. Their results show that the best performance is attained using ℓ_1 (at least theoretically). Second best is the least angle regression (LARS) algorithm. However, as shown in Table VII in [25], the LARS algorithm is quite slow. A better choice is a Tuned Two Stage Thresholding algorithm, which has good performance and is very fast. In [25], two algorithms in particular are mentioned: CoSaMP and the Subspace Pursuit algorithm. The Subspace Pursuit algorithm from [22] is shown to perform best of the two.

Recall that Θ_κ is a measurement matrix with N columns and N/κ rows and define $\mathbf{A} = \Theta_\kappa\Psi$. Then we define the Subspace Pursuit algorithm as in Algorithm 1¹. In each algorithm iteration, the pseudo-inverse is calculated as the least-squares solution as this is less computationally demanding.

Algorithm 1 Subspace Pursuit Algorithm [22]

Input:

Sparsity S , measurement and dictionary matrices combined \mathbf{A} and received, sampled signal \mathbf{y}

Initialization:

$T^0 = \{\text{indices of the } S \text{ largest absolute magnitude entries in the vector } \mathbf{A}^T\mathbf{y}\}$

$\mathbf{y}_r^0 = \mathbf{y} - \mathbf{A}_{T^0}\mathbf{A}_{T^0}^\dagger\mathbf{y}$

$\ell = 0$

repeat

$\ell \leftarrow \ell + 1$

$\tilde{T}^\ell \leftarrow T^{\ell-1} \cup \{\text{indices of the } S \text{ largest absolute magnitude entries in the vector } \mathbf{A}^T\mathbf{y}_r^{\ell-1}\}$

$T^\ell \leftarrow \{\text{indices of the } S \text{ largest absolute magnitude entries in the vector } \mathbf{A}_{\tilde{T}^\ell}^\dagger\mathbf{y}\}$

$\mathbf{y}_r^\ell \leftarrow \mathbf{y} - \mathbf{A}_{T^\ell}\mathbf{A}_{T^\ell}^\dagger\mathbf{y}$

until $\|\mathbf{y}_r^\ell\|_2 > \|\mathbf{y}_r^{\ell-1}\|_2, \ell \geq S$

To demonstrate the performance of the Subspace Pursuit algorithm with the Gold dictionary, we have performed numerical experiments to find the phase transition in the noise-less case for various choices of measurement matrices. The size of Gold dictionary used is $m = 10$, i.e. the dictionary matrix Ψ is of size 1023×1023 . The results are shown in Fig. 1. For each curve, we generate a surface plot of the rate of success, based on Monte Carlo simulations. In this surface plot, a clear transition curve is evident and to condense the results we only plot the transition curve where the probability of error crosses 0.5. Each surface plot is generated so that new simulations are conducted until the Mean Squared Error (MSE) between the i th and the $(i - 1)$ th figure is less than 10^{-5} . For each parameter set and in each simulation, an experiment is a success (1) if the MSE between the reconstructed and the received signal is less than 10^{-6} and a failure (0) otherwise. The three measurement matrices used are as follows:

- A Rademacher distributed measurement matrix, with a dense structure where entries are either -1 or 1 ,

¹In the first initialization step we choose to take the transpose of \mathbf{A} , rather than the Penrose–Moore pseudo-inverse. If instead the Penrose–Moore pseudo-inverse is used, the performance at high values of δ and ρ is increased in Fig. 1, but so is the computational complexity. This issue is not treated in more detail here, since our problems are assumed to always have low ρ .

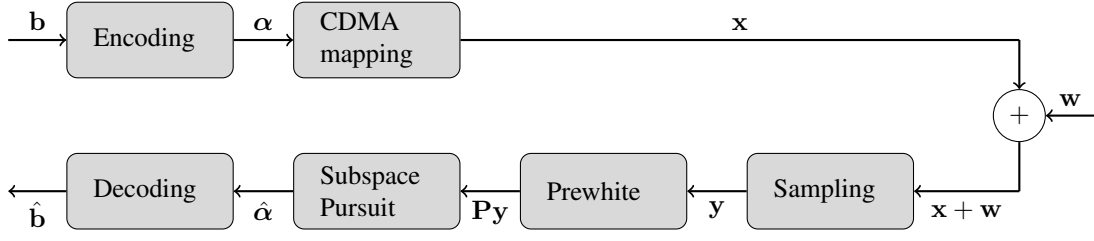


Fig. 2: Flow chart of the discrete numerical experiment.

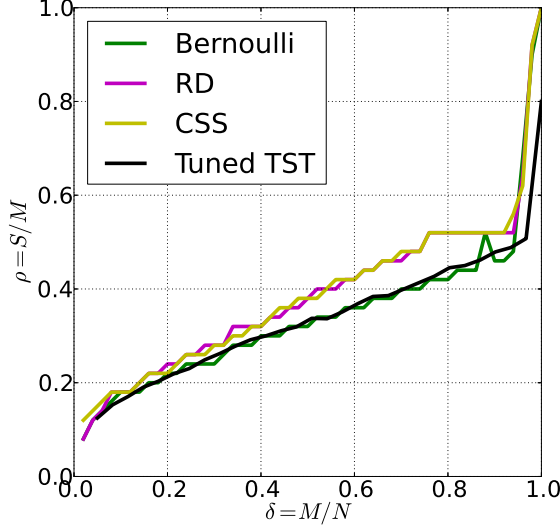


Fig. 1: Phase Transition Diagrams for the three different measurement matrices (Rademacher, RD and CSS measurement matrix) with dictionary size $m = 10$. The black line is the phase transition line for the Tuned Two Stage Thresholding (TST) algorithm from [25].

- A RD measurement matrix, with a banded structure, where entries are either -1 or 1 on the band and 0 outside, and
- the proposed CSS measurement matrix.

To validate the above results, we have inserted the phase transition line for the Tuned Two Stage Threshold (TST) algorithm from [25] in Fig. 1². As can be seen our implementation corresponds well with their results and it is clear that the proposed CSS measurement matrix performs close to identically to the RD measurement matrix and that, as previously argued, the \mathbf{D} matrix is unnecessary. Notice also the clear horizontal line in the graph around $\delta = 0.9$ and $\rho = 0.5$. We analyze this irregularity more in Section VI.

IV. DISCRETE NUMERICAL EXPERIMENT

In the above analysis, we have focused on the noise-less case and have shown that the presented dictionary and measurement matrix setup does enable CS for certain levels of sparsity. We therefore now return to the original signal model

in Eqn. (2) and investigate the noisy case by carrying out Bit Error Rate (BER) experiments. In Fig. 2 a flow chart of the numerical experiment is shown. First, we encode a randomly generated bit sequence \mathbf{b} to form the sparse vector α from Eqn. (1). The non-zero positions are chosen randomly from a uniform distribution. Each non-zero position contains a QPSK symbol. Then, α is used to create a CDMA signal using the Gold dictionary as $\mathbf{x} = \Psi\alpha$. This signal is then corrupted by additive white Gaussian noise, generated according to a chosen SNR value. Here, SNR is defined as follows:

$$\text{SNR} = \mathbb{E} \left[\frac{\|\mathbf{x}\|_2^2}{\|\mathbf{w}\|_2^2} \right] = \frac{\|\mathbf{x}\|_2^2}{N\sigma^2}, \quad (9)$$

where $w \sim \mathcal{N}(0, \sigma^2 I)$ with σ^2 the variance of the noise.

At the receiver, the sampling is modelled as in Eqn. (2) with multiplication by a measurement matrix. In the simulations we use $\kappa = 2$ or $\kappa = 4$. As is shown in the phase transition plots previously, the method also works for other choices of κ in the noise-less case. However, to clearly demonstrate that our implementation produces the expected 3 dB drop in performance per doubling of κ due to noise folding, we have chosen these two values. A measurement matrix based on samples obtained from a Rademacher distribution introduces colored noise. This decreases the performance, unless the signal is prewhitened before the reconstruction algorithm. This coloring occurs because the rows in the Rademacher matrix are not orthogonal. In the RD and CSS measurement matrices the rows are orthogonal and prewhitening is therefore unnecessary. The prewhitening is achieved by multiplying the received \mathbf{y} vector with a new matrix \mathbf{P} to obtain $\tilde{\mathbf{y}} = \mathbf{P}\mathbf{y}$. By setting $\mathbf{P} = \mathbf{C}^{-1}$, where \mathbf{C} is e.g. the Cholesky factorization ($\mathbf{C}\mathbf{C}^T = \Theta_\kappa \Theta_\kappa^T$), the variance of the noise term $\tilde{\mathbf{w}} = \mathbf{P}\Theta_\kappa \mathbf{w}$ from Eqn. (2) becomes:

$$\mathbb{E}[\mathbf{P}\Theta_\kappa \mathbf{w} \mathbf{w}^T \Theta_\kappa^T \mathbf{P}^T] = \sigma^2 \mathbf{C}^{-1} \mathbf{C} \mathbf{C}^T (\mathbf{C}^{-1})^T = \sigma^2 \mathbf{I}. \quad (10)$$

After prewhitening, we reconstruct the sparse vector $\hat{\alpha}$ using the Subspace Pursuit algorithm, which now also must include the \mathbf{P} matrix, i.e. $\mathbf{A} = \mathbf{P}\Theta_\kappa \Psi$. It is clear that this extra step increases complexity, but note that this step is only performed for the Rademacher measurement matrix. The \mathbf{A} matrix must be generated anew for each slot because a new measurement matrix Θ_κ is generated. The RD and CSS measurement matrices skip this step as their rows are orthogonal. After obtaining the sparse vector $\hat{\alpha}$, we are able to decode the original bit sequence, $\hat{\mathbf{b}}$.

To validate the obtained results, we compare the numerical results with the theoretical performance for non-coherent

²Data from <http://sparselab.stanford.edu/OptimalTuning/main.htm>

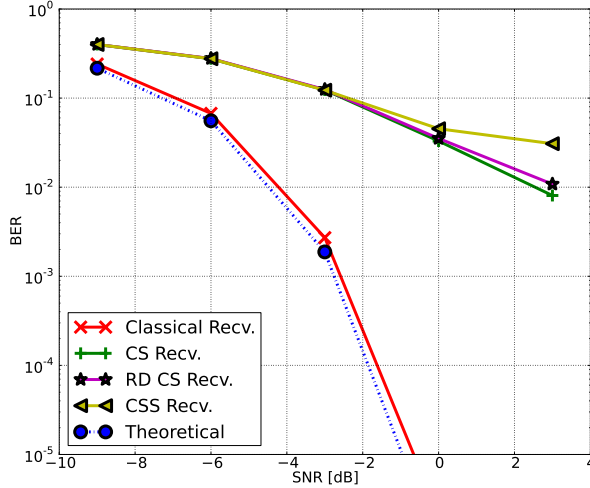
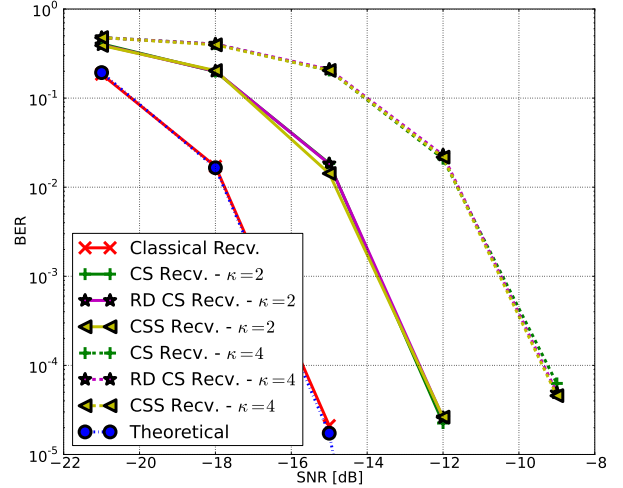
(a) $m = 5, \kappa = 2$.(b) $m = 10$, solid lines are for $\kappa = 2$, dashed are $\kappa = 4$.

Fig. 3: BER versus SNR for different dictionary sizes and choices of κ . CS here is the Rademacher measurement scheme. Simulations were run until 100 bit errors were found for each SNR point.

MFSK [26]:

$$P_b = \frac{N}{2(N-1)} \frac{1}{N} \sum_{k=2}^N (-1)^k \binom{N}{k} \exp \left(N \cdot \text{SNR} \left(\frac{1}{k} - 1 \right) \right), \quad (11)$$

where SNR is the signal to noise ratio. We use the non-coherent formula because the CDMA codes are QPSK modulated. This corresponds to a phase shift of the original codes, which makes the receiver non-coherent. Furthermore, for the above result to hold, we must fix $S = 1$, i.e. the CDMA signal is 1-sparse. Then, instead of performing reconstruction of the sparse α , we may instead perform classification as in [13], [14]. This would replace the Subspace Pursuit algorithm with a simpler estimation framework. However, to conserve generality and because we use $S = 10$ later, we continue using the general CS framework and the Subspace Pursuit algorithm.

As dictionary we use Gold sequences with $m \in \{5, 10\}$. This reveals the performance for different dictionary sizes and especially $m = 10$ is interesting as it is the LFSR length used in e.g. GPS. The results of the numerical experiments are shown in Fig. 3.

As can be seen, for $m = 5$ both the Rademacher, RD and especially the CSS measurement matrix seems to perform poorly. For high SNR values there is more than the expected 3 dB loss per octave due to noise folding. At $m = 10$ the loss is almost exactly 3 dB per halving of the sampling rate. For $m = 10$ we have also included the result for $\kappa = 4$ to show the performance when the sampling rate is reduced to a quarter of the Nyquist sampling rate. Again, the curve follows the previous results for noise folding, as the performance degrades by approximately 3 dB more for all the CS-enabled receiver structures. These results show

that the CSS measurement matrix, though simpler than all the other measurement matrices, performs equally well in the above experiments for $m = 10$. For small dictionary sizes, its performance is worse.

V. RF NUMERICAL EXPERIMENT

To obtain more realistic communication-relevant results, we have extended the above discrete numerical experiment to a full transmitter/receiver simulation with RF up and down-conversion and with root raised cosine pulse shaping and matched filter. This we have done to demonstrate that the results from Fig. 3 translate to a realistic transmitter/receiver system. The construction of the experiment is visualized in Fig. 4. This conceptual flow chart also visualizes how the ADC process must be incorporated in a receiver structure to implement the proposed CSS method. The experiment we have conducted is based on a QPSK signal with a chip rate of 10^6 chips per second using a root raised cosine pulse shaping filter with a roll-off factor of 1. This signal is represented in the simulation as sampled at 10 times that rate, to emulate an analog signal. The signal is up-converted to an RF frequency of 3 MHz, i.e. 3 times the chip rate. The RF signal is sampled at 12 MHz, again to emulate an analog signal. Here, AWGN is added followed by down-conversion again. The down-conversion is implemented as perfect direct down-conversion. This is accomplished by first multiplying with a complex exponential, followed by taking an FFT of the signal. In the output from the FFT, all values above the chip rate are set to 0, after which the inverse FFT is taken. At baseband, the sampling is done by a matched filter based on the same root raised cosine that is used for pulse shaping. The samples are then input to the Subspace Pursuit algorithm, similar to the discrete numerical experiment.

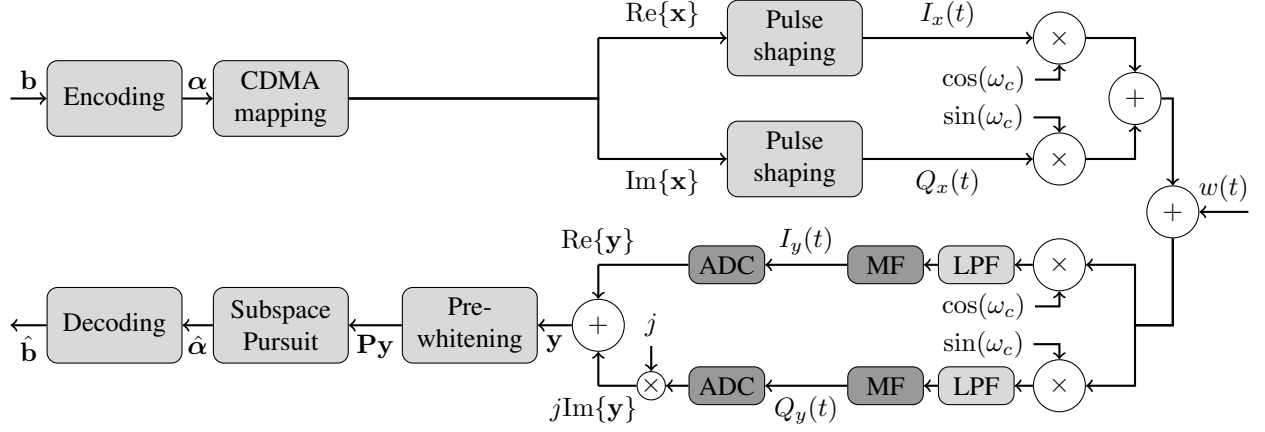
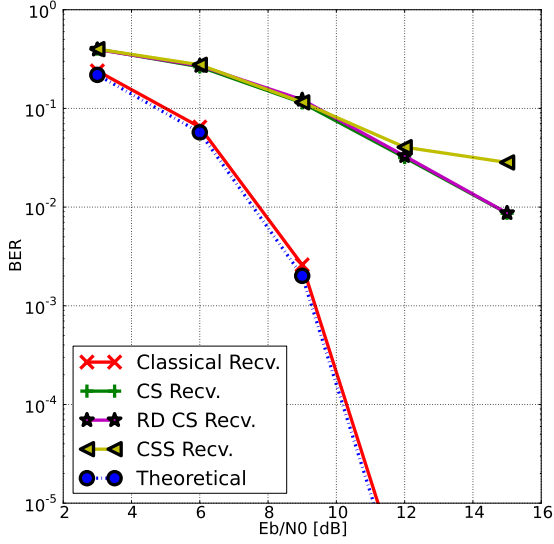
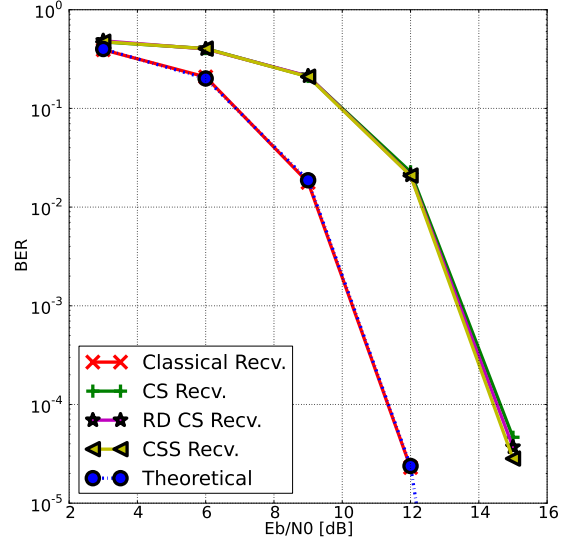


Fig. 4: Conceptual flow chart of the RF numerical experiment. Note that all continuous variables here are only conceptual. In the numerical experiments they are represented as discrete, oversampled sequences. Here, MF is a matched filter and LPF is a low-pass filter. Dark boxes signify components that must be changed compared to a traditional architecture to enable the CS subsampling described in this work.



(a) $m = 5$, 100 errors found for each point, $\kappa = 2$.



(b) $m = 10$, 100 errors found for each point, $\kappa = 2$.

Fig. 5: BER versus E_b/N_0 for different dictionary sizes. CS here is the Rademacher measurement scheme. Simulations were run until 100 bit errors were found for each E_b/N_0 point.

The results of the experiment are shown in Fig. 5. The theoretical curve is calculated using a modified version of the non-coherent MFSK equation used before:

$$P_b = \frac{M}{2(M-1)} \frac{1}{M} \sum_{k=2}^M (-1)^k \binom{M}{k} \exp \left(\log_2(4) \frac{E_b}{N_0} \left(\frac{1}{k} - 1 \right) \right), \quad (12)$$

where E_b/N_0 is the energy per bit per noise spectral density and we multiply E_b/N_0 with $\log_2(4)$ because there are 4 constellation points in QPSK. As can be seen, the results here are close to identical with those for the simpler discrete

numerical experiment. Noise folding still gives rise to a penalty, which makes CS a trade-off between sampling rate and BER performance. However, previous work has suggested that quantization may shift the trade-off point, so that CS obtains both the low sampling rate and a better performance than a classical receiver [18]. We investigate this in the following.

A. RF Numerical Experiment with Quantization

In [18], it is proposed to combat noise folding with quantization as a CS receiver is able to quantize the sampled signal better, since it takes fewer measurements. By better quantization we mean that if the CS receiver takes half as many

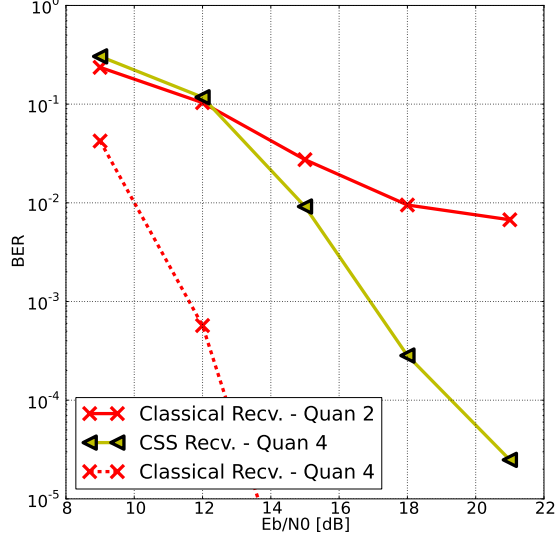


Fig. 6: BER versus E_b/N_0 for a classical receiver and a CSS receiver, both with quantization. $m = 7$, $S = 10$, $\kappa = 2$ and 100 errors found for each E_b/N_0 point. The dotted line is for a classical receiver with 4 bits of quantization per sample.

samples, it may quantize twice as well at no additional cost. We have investigated this by applying uniform quantization to the RF experiment performed in the previous section. However, as simple QPSK modulation is used, only the sign matters for demodulation and therefore quantization has no effect in the simple case of $S = 1$ used so far. Therefore, we investigate $S = 10$ instead and used 2 bits of quantization per sample (i.e. 4 bits of quantization for CSS as $\kappa = 2$). This is merely intended as an example study to show that when taking into account quantization, CS may perform better than a classical receiver. The result of the numerical experiment is shown in Fig. 6. As can be seen, quantization makes CS a better alternative in this scenario. The CS approach becomes significantly better for high E_b/N_0 values, because the classical receiver is not able to quantize the signal properly. For comparison, we have also included the same result for a classical receiver with 4 bits of quantization, i.e. the same level of quantization as the CSS receiver. Then it becomes clear that the classical receiver again is the best choice, but remember that it operates at twice the sampling frequency. A CS-enabled receiver can therefore be seen as a trade-off point between sampling rate and dynamic range.

VI. COMPLEXITY ANALYSIS

To evaluate the Subspace Pursuit algorithm, we investigate the computational complexity of the algorithm, shown in Tab. I, where K is the number of iterations used in the Subspace Pursuit algorithm, S is the sparsity, M is the number of measurements taken and N is the number of Nyquist samples.

The matrix \mathbf{A} is real, but since \mathbf{y} is complex this affects the matrix-vector computations. A matrix-vector product then

costs $4MN$ and calculating a residual costs $2M + 8MS$.

The pseudo-inverse is never calculated, instead a linear least-squares problem is solved using the Singular Value Decomposition (SVD). Solving a least-squares problem with S variables and M observations using the SVD costs [27]:

$$\text{Cost}_{\text{LS with SVD}} \sim 2MS^2 + 11S^3. \quad (13)$$

Notice that the first least square problem in the loop takes in $2S$ atoms from the dictionary. The cost of sorting and locating entries is not taken into account here, as those algorithms are more memory then computationally demanding.

It is also important to notice that the problem sizes involved here are very small. Compressive sensing only works for sparse signals, so S is often small compared to M and N . In the examples given here, $N = 1023$ is the largest dimension we have worked with. Because of this, the mathematical model in Tab. I is not adequate, as the computational complexity is instead dominated by programming language overhead, such as the cost of calling different functions. Therefore, it is important to include an extra term: cK , where K is the number of iterations performed and c is some constant that depend on system and programming language overhead.

It is of interest to investigate the required number of iterations, K , of the Subspace Pursuit algorithm, to better understand the cost of using CSS. In Fig. 7 we show the number of iterations used to generate the results in Fig. 1. The horizontal line through $\rho = 0.5$ is interesting and unexpected. If we change the input sparsity to the Subspace Pursuit algorithm from S to $2S$, the line moves from $\rho = 0.5$ to $\rho = 0.25$, which means it is related to the number of atoms available to the Subspace Pursuit algorithm in each iteration. It is not related to the dictionary type, as we have obtained exactly the same phase transition diagrams and iteration counts with a Haar wavelet packet dictionary. Furthermore, it is not due to a "lucky" initial guess, as the line first emerges in the

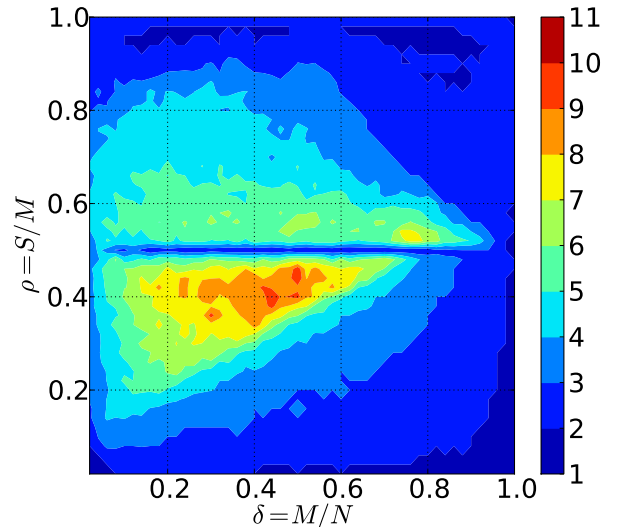


Fig. 7: Number of Subspace Pursuit iterations for the CSS measurement matrix and Gold Dictionary size $m = 10$.

TABLE I: Computational cost of the Subspace Pursuit algorithm.

Action	Approx. cost
<i>Initialization:</i>	
• 1 computation of $\mathbf{A}^T \mathbf{y}$	$4MN$
• 1 computation of $\mathbf{y} - \mathbf{A}_{T^0} \mathbf{A}_{T^0}^T \mathbf{y}$	$2M + 8MS$
<i>Loop:</i>	
• K computations of $\mathbf{A}^T \mathbf{y}$	$4KMN$
• K least squares problems $(\mathbf{A}_T^\dagger \mathbf{y})$	$K(2M(2S)^2 + 11(2S)^3)$
• K computations of $\mathbf{y} - \mathbf{A}_{T^0} \mathbf{A}_{T^0}^T \mathbf{y}$	$K(2M + 4MS + 2MS^2 + 11S^3)$
Total:	$99KS^3 + 4(K+1)MN$ $+ 2(K+1)M + 4(K+2)MS + 10MK S^2$

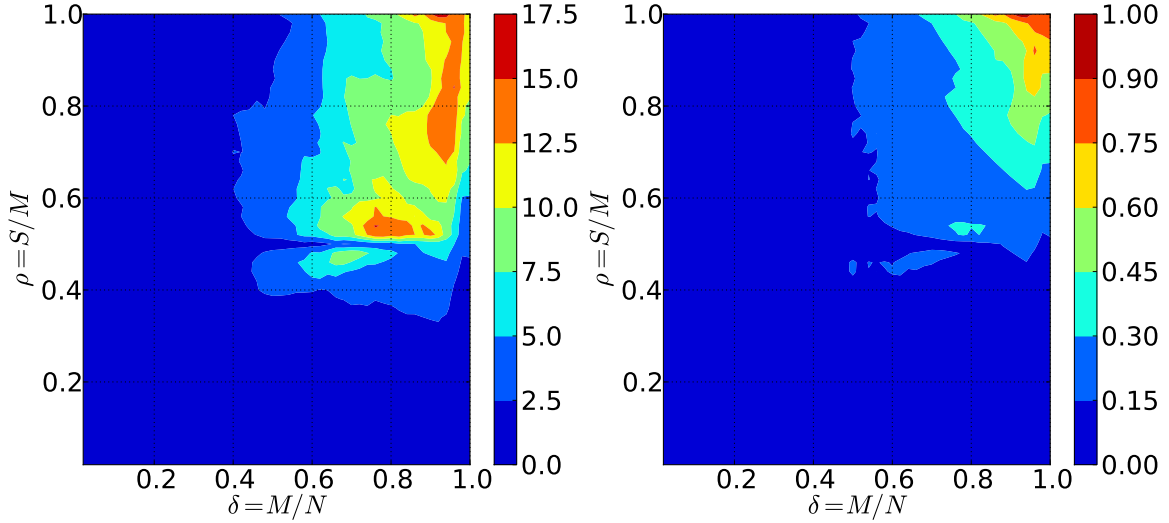


Fig. 8: Computational cost of the Subspace Pursuit algorithm for the CSS measurement matrix and Gold Dictionary size $m = 10$. The figure to the left shows measured execution time in seconds from the numerical experiment conducted in Fig. 1 and the figure to the right is generated using the formula for the total computational cost found in Tab. I plus the term cK with $c = 3 \cdot 10^9$. For the figure to the right, the number of iterations of the Subspace Pursuit algorithm, K , is taken from Fig. 7 and the numbers are normalized.

third iteration of the algorithm. It seems to be an overlooked property of the algorithm, which has gone unnoticed so far because the line in Fig. 7 lies in the region of Fig. 1, where the algorithm cannot find the correct solution anyway.

Finally, we have measured the computation time for running the Subspace Pursuit algorithm for the CSS numerical experiment in Fig. 1. These are compared to the theoretical values obtained by using Tab. I. The constant c has been set to $3 \cdot 10^9$, which is a value found to give a good accordance with the numerically found values. It is important to note that this choice of c is very much a function of the algorithm, problem size, programming language and the machine on which the experiment is conducted and should therefore not be seen as a general choice. The result is shown in Fig. 8. The values in the figure on the right are normalized to one, as they are completely dependent on machine power and are only shown here to visualize how much the computational requirements change with the parameters. As can be seen, the numerically obtained computation times seem to correspond fairly well to the mathematical model. Each point in the above numerical

experiment has been run as a simulation on 1 out of 16 threads on computation nodes with 2x Intel Xeon X5570 CPUs and 48GB memory.

VII. CONCLUSION

In this work we apply CS to a general CDMA system and we show that it is possible to use a very simple measurement scheme at the receiver side to enable subsampling of the CDMA signal. We show that the performance of the proposed receiver scheme is affected negatively in BER performance, similar to other CS schemes. However, we also show that when taking quantization into account, the proposed receiver model performs better in our example than a classical receiver with the same quantized bit rate. Finally, we investigate the complexity of the developed algorithms and compare the computational cost of the numerical experiments with the theoretically calculated computation cost.

Our work here has shown that CS used in spread spectrum receivers allows for a simplified front-end compared to other

state-of-the-art CS sampling designs. Furthermore, we have shown that the problem of noise folding may be remedied in some cases by using quantization. Future work should investigate further which scenarios may benefit from CS and also perform laboratory experiments with the CSS receiver structure. Furthermore, the premise of this work is that taking fewer samples conserves power. This must be validated through laboratory experiments and the power efficiency of the CSS receiver structure should be better evaluated.

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